

On the Uniform Computational Content of the Baire Category Theorem

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joint work with

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Baire category theorem

Let X be a complete metric space.

Theorem (Baire category theorem)

Let $(A_n)_{n \in \mathbb{N}}$ be closed nowhere dense subsets of X .

$$\bigcup_{n \in \mathbb{N}} A_n \subsetneq X$$

Representation of closed sets

- **negative information**, $\mathcal{A}_-(X)$, ϕ_- ,
 A is the complement of an open set given as an enumeration of basic open balls
- **positive information**, $\mathcal{A}_+(X)$, ϕ_+ ,
 A is the closure of an enumeration of points in X .
- **set of cluster points**, $\mathcal{A}_*(X)$, ϕ_* ,
 A is the set of cluster points of an enumeration of points in X .

Example: Closed Choice

Let X metric space.

Definition

$$C_X : \subseteq \mathcal{A}_{\square}(X) \rightrightarrows X \\ A \mapsto A$$

- **positive information** ($\square = +$),
 C_X is trivial.
- **set of cluster points** ($\square = *$),
 C_X is the same as finding a cluster point (similar as Bolzano-Weierstraß theorem)
- **negative information** ($\square = -$),
right formulation, non-continuous.

Baire category theorem

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Theorem (Baire category theorem)

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Formulate as computational problem:

BCT₀ Given $(A_n)_{n \in \mathbb{N}}$ closed nowhere dense. There is an $x \in X \setminus \bigcup_{n \in \mathbb{N}} A_n$.

$$\text{BCT}_0 : \subseteq \mathcal{A}_-(X)^{\mathbb{N}} \Rightarrow X$$

BCT₁ Given $(A_n)_{n \in \mathbb{N}}$ closed, such that $\bigcup_{n \in \mathbb{N}} A_n = X$. There is an index n such that A_n is somewhere dense.

$$\text{BCT}_1 : \subseteq \mathcal{A}_-(X)^{\mathbb{N}} \Rightarrow \mathbb{N}$$

BCT₂, **BCT₃** are defined like **BCT₀** and **BCT₁** but with positive input.

Overview over the strength of BCT

		classical reverse mathematics
BCT_0	computable	RCA_0
BCT_1	equivalent to $C_{\mathbb{N}}$ implies Banach inverse mapping theorem, etc.	RCA_0
BCT_2	computability theoretic version related to 1-generic, forcing, $BCT//$	Π_1^0G
BCT_3	equivalent to cluster point problem Space X has to be perfect (no isolated points.) E.g., $2^{\mathbb{N}}$, $\mathbb{N}^{\mathbb{N}}$.	Π_1^0G

In classical reverse mathematics (Brown, Simpson '93)
 $RCA_0 + BCT// \vdash$ "Banach inverse mapping theorem".

BCT on non-perfect space X

If X has isolated points.

Proposition

$$\text{BCT}_2 \equiv_{\text{sW}} id_0$$

$$\text{BCT}_3 \equiv_{\text{sW}} id_{\mathbb{N}}$$

In particular, BCT_2 , BCT_3 are computable in this case.

BCT on perfect spaces

Theorem (Brattka, Hendtlass, K.)

BCT_i for a perfect polish space X is sW-equivalent to BCT_i for $\mathbb{N}^{\mathbb{N}}$.

Proof.

- BCT_i for X is reducible to BCT_i for $\mathbb{N}^{\mathbb{N}}$:
 - Let $\delta: \mathbb{N}^{\mathbb{N}} \rightarrow X$ be the Cauchy representation.
 - Take $\delta BCT_i(\delta^{-1}(A_n))$.
- BCT_i for $2^{\mathbb{N}}$ is reducible to BCT_i for X :
 - By perfectness there is an embedding $\iota: 2^{\mathbb{N}} \rightarrow X$ with computable inverse. (Brattka, Gheradi '08)
 - $A \subseteq 2^{\mathbb{N}}$ nowhere dense $\Rightarrow \iota(A)$ nowhere dense.
 - Take $\iota^{-1}(BCT_i((\iota(A_n))_n))$.
- BCT_i for $\mathbb{N}^{\mathbb{N}}$ is reducible to BCT_i for $2^{\mathbb{N}}$:
 - Embed $\mathbb{N}^{\mathbb{N}}$ into $2^{\mathbb{N}}$ via $p \mapsto 1^{p(0)}01^{p(1)}01^{p(2)}0\dots$.
 - Range is (c.e.) comeager. □

Consider now only $X = \mathbb{N}^{\mathbb{N}}$.

Theorem (Brattka '01, Brattka, Gherardi '11)

- $C_{\mathbb{N}} \equiv_{sW} BCT_1$,
- $CL_{\mathbb{N}} \equiv_{sW} BCT_3 \equiv_{sW} BCT'_1$.

Theorem (Brattka, Hendtlass, K.)

- BCT_0, BCT_2 are densely realized.
In particular $C_2 \not\equiv_W BCT_0, BCT_2$.
- $BCT'_0 \equiv_{sW} BCT_2$.

Proof of $BCT_2 \equiv_{sW} BCT'_0$ and $BCT_3 \equiv_{sW} BCT'_1$

Proposition

$id_{+-} : \mathcal{A}_+(X) \rightarrow \mathcal{A}_-(X) \leq_{sW} \text{lim}$

Gives $BCT_2 = BCT_0 \circ id_{+-} \leq_{sW} BCT'_0$ and $BCT_3 \leq_{sW} BCT'_1$.

Proposition (Brattka, Gherardi, Marcone '12)

$id : \mathcal{A}_*(X) \rightarrow \mathcal{A}_-(X)'$ is a computable isomorphism.

Proposition

There is an $M : \subseteq \mathcal{A}_*(X) \rightrightarrows \mathcal{A}_+(X)$ such that,

- $M(A) \subseteq \{B : A \subseteq B\}$
- A nowhere dense $\Rightarrow B \in M(A)$ nowhere dense. (X perfect)

The mapping

$$\mathcal{A}_-(X)' \xrightarrow{id} \mathcal{A}_*(X) \xrightarrow{M} \mathcal{A}_+(X)$$

gives $BCT'_0 \leq_{sW} BCT_2$ and $BCT'_1 \leq_{sW} BCT_3$.

1-generic

A point $p \in 2^{\mathbb{N}}$ is 1-generic **relative to** q if it meets or avoids any c.e. open set U_i^q , i.e.,

$$\exists w \sqsubseteq p \left(w2^{\mathbb{N}} \subseteq U_i^q \quad \text{or} \quad w2^{\mathbb{N}} \cap U_i^q = \emptyset \right).$$

Equivalently: $p \notin \partial U_i^q$

Theorem

$$\text{BCT}_0 \leq_{\text{sW}} \text{1-GEN} \leq_{\text{sW}} \text{BCT}_2$$

Proof.

- 1 For nowhere dense A , $A = \partial A = \partial A^c$.

$$\text{BCT}_0((A_i)_i) = 2^{\mathbb{N}} \setminus \bigcup_{i=0}^{\infty} A_i = \bigcap_{i=0}^{\infty} (2^{\mathbb{N}} \setminus \partial A_i^c)$$

Now $A_i^c = U_j^q$ for a suitable j . Thus, $\text{BCT}_0 \leq_{\text{sW}} \text{1-GEN}$.

- 2 Use $\text{BCT}_2 \equiv_{\text{sW}} \text{BCT}'_0$ and compute ∂U_i^q in the limit. □

1-generic (cont.)

Theorem

$BCT_0 <_{sW} 1\text{-GEN} <_{sW} BCT_2$
(The implications are strict.)

Proof sketch.

- 1 Sufficient to use a **weakly** 1-generic in the previous proof.
Apply the fact that there are weakly 1-generics that are not 1-generic.
 $BCT_0 \leq_{sW} 1\text{-weakGEN} <_{sW} 1\text{-GEN}$.

- 2 (Uniform) Theorem of Kurtz shows that $1\text{-GEN} \leq_{sW} WWKL'$.^a

Lemma of Kučera shows that $WWKL'$ can be realized such that its output is low for Ω .

There is a computable p such that $BCT_2(p)$ is not low for Ω .

Thus, $BCT_2 \not\leq_W WWKL'$. □

^aActually, $(1 - *)\text{-}WWKL'$

Definition ($\Pi_1^0 G$, classical reverse math)

Let $D_i \subseteq 2^{<\mathbb{N}}$ be a sequence of dense, uniformly Π_1^0 -set.

There is a set $G \subseteq 2^{\mathbb{N}}$ meeting each D_i , i.e., $\exists s \in D_i (s \sqsubseteq G)$.

$\Pi_1^0 G$ related to forcing constructions.

Formulation in the Weihrauch lattice: Model properties of D_i using a suitable representation.

Definition

$\phi_{\#}(p) = A : \iff \phi_{-}(p) = D$ and $A = 2^{\mathbb{N}} \setminus \bigcup_{w \in D} w 2^{\mathbb{N}}$,
where $D \subseteq 2^{<\mathbb{N}}$.

Definition ($\Pi_1^0 G$, Weihrauch version)

$$\Pi_1^0 G := \subseteq \mathcal{A}_{\#}(2^{\mathbb{N}})^{\mathbb{N}} \rightrightarrows 2^{\mathbb{N}} \quad (A_n)_n \mapsto \bigcap 2^{\mathbb{N}} \setminus A_n,$$

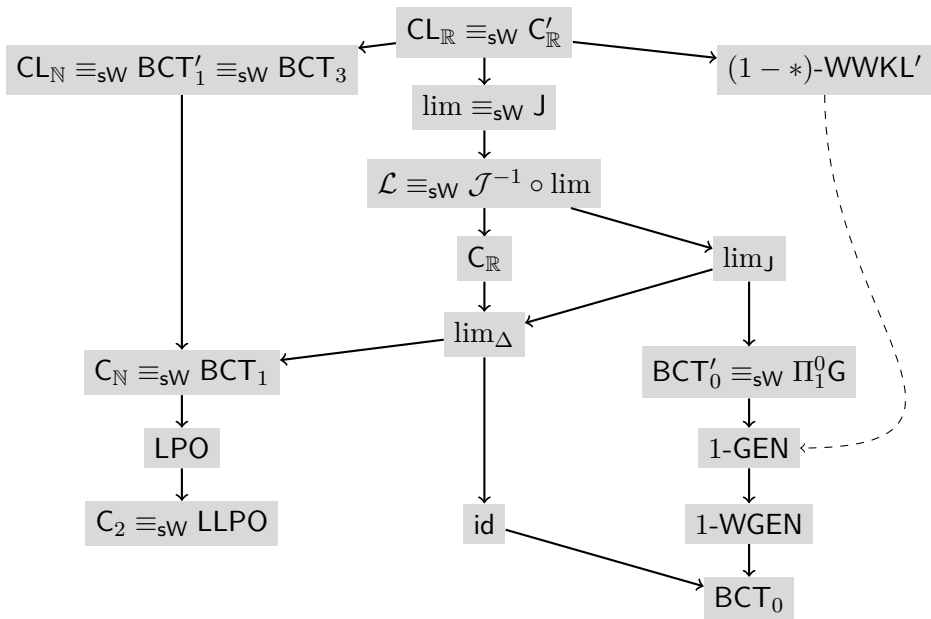
with $\text{dom}(\Pi_1^0 G) := \{(A_i)_i \mid A_i^{\circ} = \emptyset\}$.

Proposition

$\text{id}: \mathcal{A}_-(2^{\mathbb{N}})' \rightarrow \mathcal{A}_{\#}(2^{\mathbb{N}})$ is a computable isomorphism.

Corollary

$\Pi_1^0 G \equiv_{sW} \text{BCT}'_0 \equiv_{sW} \text{BCT}_2$



Summary

- Different ways of writing Baire category theorem as forall-exists statement are natural.
- Basically 2 different variants.
- Calculus characterization of $\Pi_1^0 G$.

Thank you for your attention!

Bibliography



V. Brattka, M. Hendtlass, A. Kreuzer

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[arXiv:1501.00433](https://arxiv.org/abs/1501.00433)